

5. Real-Time Programming

5.2 Real-Time Java

Roadmap of Section 5.2

- Real Time Specification for Java
- RTSJ Features
 - RealtimeThreads
 - Memory Management
 - RawMemoryAccess
 - Asynchronous Transfer of Control
 - Asynchronous Event Handling
- Reference Implementation / Available Impl.
 - Status of RTSJ
- J Consortium

History

- Dec. 1998 java specification request for real time extension for java
- Expert group – SUN, IBM, QNX Software Lab, Nortel, Rockwell, Timesys ..
- Greg Borella (IBM) first specification lead
- Sept. 1999 first public review of specification
- Late in 2001 Timesys volunteered to create the reference implementation
- Final Specification 12/11/2001
- 2003 Sun announced Mackinac project: first commercial implementation of RTSJ

Motivation

- Usage of advantages of Java
 - Cross-platform capabilities
 - Object orientation, Type Safety
 - Developers and Tools available, Rapid Application Development
- Improve real-time properties of java
 - Deterministic execution times
 - Specify real-time scheduling / known start / stop times of threads
 - Specify sufficient memory management
 - Direct access to hardware / memory

Real Time Specification for Java (RTSJ)

Java Architecture

| | |
|-----------------|----------------|
| Java Sourcecode | Java Libraries |
| JVM | |
| OS | |

Real-time Java

| | |
|-----------------|------------------------------|
| Java Sourcecode | Java+ JavaRealTime Libraries |
| RTJVM | |
| RTOS | |

- Standard Java API + Real-time Extensions :
- `javax.realtime.*`

RTSJ: Major Specification Features

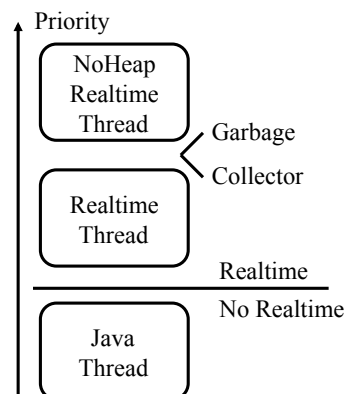
- Real-time threads with precise defined scheduling
- Mechanisms that support writing code that is not influenced by garbage collection
- Asynchronous event handlers to handle events from outside the virtual machine
- Asynchronous transfer of control
- Mechanisms that allow to control where objects will be allocated in memory
- Direct memory access

RTSJ Scheduling

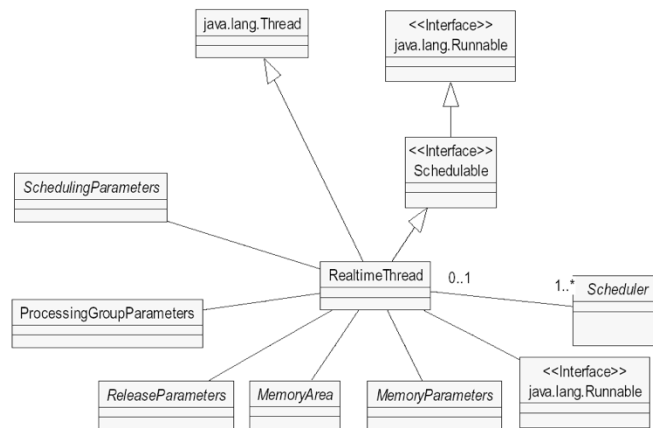
- Scheduling manages scheduling / dispatching of schedulable objects
- Schedulable object – implements *Schedulable*
- RTSJ specifies default scheduling algorithm
 - Fixed-priority preemptive scheduling
 - FIFO
 - At least 28 scheduling priorities
 - Highest priority thread always runs
- Custom scheduler can be implemented

Threads

- NoHeapRealtime Threads
 - Hard real-time
 - Higher priority than gc
 - No references to heap memory
- Realtime Thread
 - Soft real-time
 - Can be interrupted by gc
 - References to heap allowed



RealTimeThread



Periodic Threads

```

int pri = PriorityScheduler.instance().getMinPriority()+10;
PriorityParameters prip = new PriorityParameters(pri);
RelativeTime period = new RelativeTime(20 /* ms */, 0 /* ns */);
PeriodicParameters perp = new PeriodicParameters
    (null, period, null, null, null, null);
RealtimeThread rt= new RealtimeThread(prip, perp) {

    public void run() {
        int n=1;
        while (waitForNextPeriod() && (n<100)) {
            System.out.println("Hello "+n);
            n++;
        }
    }
};
rt.start();
    
```

Scheduler

- Default Scheduler : *PriorityScheduler*
 - No change of priority during runtime
- Performs feasibility analysis for sets of schedulable objects
- Cost overrun handler / missed deadline handler per process
- Controlled via *SchedulingParameter*
- Additional Scheduler must implement abstract base class *Scheduler*
- Can be installed via : *RealtimeThread*.

```
public void setScheduler(Scheduler scheduler)
```

T_1, T_2, \dots, T_n - Tasks to be performed in the real time system

C_1, C_2, \dots, C_n - Cost of each task (how long it takes to run each task)

R_1, R_2, \dots, R_n - Release time for each task (time that task becomes available to run)

D_1, D_2, \dots, D_n - Deadline for each task (when each task needs to be complete)

Asynchronous Event Handling

- Real-time and embedded systems are coupled to the real world
- Events in the real world are asynchronous
- RTSJ specifies a mechanism to bind a schedulable object to the occurrence of an event
- When the event occurs the object's run state changes to ready-to-run and is scheduled according its parameters
- Implementation should support hundreds of ev.

Asynchronous Event Handling

- An instance of `AsyncEvent` represents something that can happen
- `AsyncEventHandler` implements `Schedulable`
 - `RealTimeThread` / `NoHeapRTThread`
- Default Constructor : All properties inherited from current thread
- An instance of `AsyncEventHandler` has a method `handleAsyncEvent()` which contains the logic that should execute when the event occurs
- Method `run()` invokes `handleAsyncEvent()`

AsynchEvent Class

- `public synchronized void addHandler
 (AsyncynchronousEventHandler handler)`
 - Adds a handler to the set defined for this `AsynchEvent`
- `public void bindTo(String happening)`
 - Binds this `AsynchEvent` to an external event (a happening)
 - Happening is an implementation dependent value that binds this `AsynchEvent` to some external event
- `public synchronized void fire()`
 - Schedules the `run()` method of each handler associated with this event

Interrupt Handling Example

```
import java.realtime.*;
public class HardwareInterruptExample extends AsyncEvent{
    private int interruptNum;
    public HardwareEventExample(int num) {
        interruptNum = num;
    }
    public void setHandler(AsyncEventHandler h) {
        super.addHandler(h);
        super.bindTo(interruptNum);
    }
    class HardwareEventHandler extends AsyncEventHandler{
        private int interruptCount = 0;
        public void handleAsyncEvent(){
            interruptCount++;
            // Driver code follows
        }
    }
}
```

Time

- „Allow *description of a point in time* with up to *nanosecond accuracy and precision* (actual accuracy and precision is dependent on the precision of the underlying system).“
- „Allow *distinctions between absolute points in time*, times relative to some starting point, and a new construct, *rational time*, which allows the efficient expression of occurrences per some *interval of relative time*.“
- Abstract HighResolutionTime implements Comparable
- RelativeTime, AbsoluteTime, RationalTime

Timers

- Triggers behaviour at a particular point in time
- Special form of asynchronous events
- *OneShotTimer*
 - Fires off once at the specified time
- *PeriodicTimer*
 - Fires off at the specified time and then
 - periodically with a specified interval
- Clock : interface to the system's real-time clock

Timer Example

```
PeriodicTimer pt = new PeriodicTimer(  
    new RelativeTime(200,0),  
    new RelativeTime(200,0),null);  
  
ReleaseParameters rp = pt.createReleaseParameters();  
  
pt.addHandler(new AsyncEventHandler  
    (null,rp,null,null,null) {  
    public void handleAsyncEvent() {  
        System.out.println("Timer went off ");  
    }  
});  
pt.start(); // start the timer
```

Asynchronous Transfer of Control

- Allows interrupting a thread by raising interrupted exceptions
- One thread can throw an exception into another thread
- Better way of notifying the application about the occurrence of a significant event
- Behaves like `Thread.stop(deprecated)` but is safer
- Can be used as a time-out mechanism
- Asynchronous exception deferred if thread is in synchronized block or uninterruptible method
 - Methods can be made interruptible if *AsynchronouslyInterruptedException* is added to throw clause

Asynchronously Interrupted Exception

- A thread that wants to be interrupted when significant events occur, should mark its methods as throwing *AsynchronouslyInterruptedException*
- The thread would not be interrupted if it is executing a method that is not marked as throwing *AsynchronouslyInterruptedException*
- Triggered when *RealtimeThread.interrupt()* is called

Memory Management

- **Definition of memory areas for object allocation**
- **Heap memory – no real-time**
 - Standard Java Heap (one per Virtual Machine)
- **Immortal memory – real-time capable**
 - Allocated objects exist until the end of the application
- **Scoped memory – real-time capable**
 - Manual memory management (defined scope)
- **Physical memory areas**

Scoped Memory

- **Activated using the method enter**
 - `public void enter(Runnable r)`
- **All allocation in run-method of runnable are done in ScopedMemory**
- **All objects in Scoped memory will be finalized and collected if :**
 - Last real-time thread referencing the scoped exits
- **Reference counting of real-time thread using the scope**
- **Single Parent rule for Scope Stacks**
 - No cycles in scope dependencies

Memory Management

Scoped Memory - Types

- **VTMemory**

- Allocation may take a variable amount of time
- Not subject to garbage collection

- **LTMemory**

- Not subject to garbage collection
- Guarantees linear execution time for object allocations from the area

- **(CTMemory) in jRate**

- Allocation in constant time

ScopedMemory Example

```
Final ScopedMemory myScope = new VTMemory();
myScope.enter(new Runnable()
{
    public void run()
    {
        ...           // all new calls here are
                       // allocated to myScope
    }
}                   // end of run
)                   // end of enter
```


ScopedMemory Example 2

```
final ScopedMemory s = new LTMemory (16,1024);

RealtimeThread t = new RealtimeThread (null,
                                         null,
                                         new MemoryParameter (s),
                                         null,
                                         new Runnable ()
                                         {
                                             public void run ()
                                             {
                                                 // ...
                                             }
                                         }
                                         );
```

Nested Scoped Memory

```
Runnable nestedLogic = new Runnable() {
    public void run() {
        MemoryArea ma2 = new LTMemory(...);
        Runnable moreNestedLogic = new Runnable() {
            public void run() {A a = new A();}
            ma2.enter(moreNestedLogic);
        };
    };
};

MemoryArea ma1 = ...
ma1.enter(nestedLogic);
```

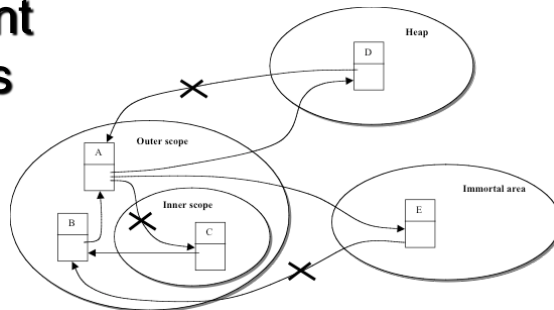
Immortal Memory

- Shared among all threads
- Objects allocated within `ImmortalMemory` live until the end of the application
 - Objects still exist without any reference to it
- Can be scanned by garbage collector, but not collected itself
- Singleton class
- `ImmortalPhysicalMemory`

Budgeted allocation

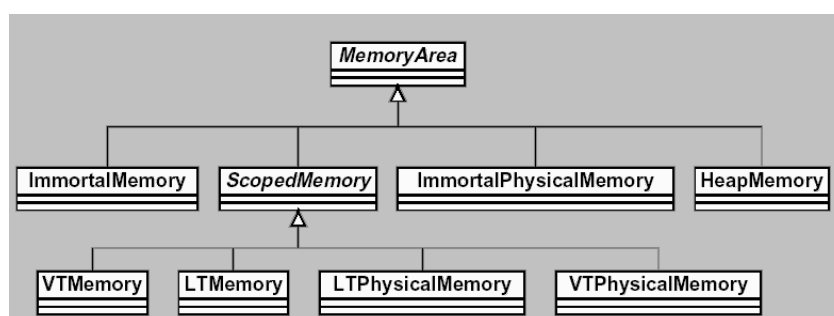
- RTJS provides limited support for memory allocation budgets
- Maximum memory area consumption and maximum allocation rates for real-time threads
- Definition in `MemoryParameter` of `RealTimeThread` constructor

Assignment restrictions



| | Reference to Heap | Reference to Immortal | Reference to Scoped |
|----------------|-------------------|-----------------------|--------------------------------------|
| Heap | Yes | Yes | No |
| Immortal | Yes | Yes | No |
| Scoped | Yes | Yes | Yes, if same, outer, or shared scope |
| Local Variable | Yes | Yes | Yes, if same, outer, or shared scope |

Memory Area - Classes



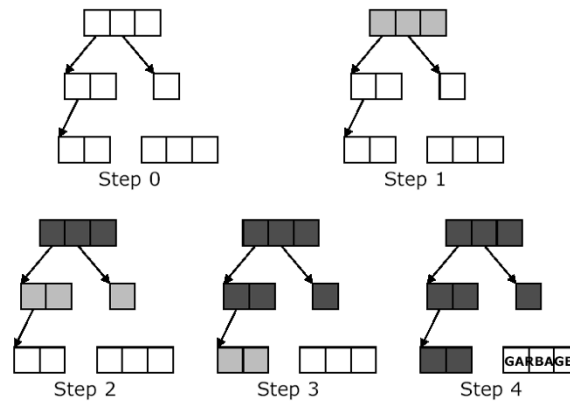
Garbage Collection

- **Reference Counting**
- **Mark-and-Sweep**
 - Distinguish live objects from garbage
 - Start in local variable array, operand stack
 - Mark all referenced objects alive
 - Remove all unmarked objects
- **Mark-and-Compact**
 - Adds de-fragmentation to mark-and-sweep algorithm

Real-time garbage collection

- **Fine-grained incremental garbage collection**
- **Garbage collection should run interleaved with normal threads – not atomic !**
- **Incremental tracing collectors**
 - Objects traversed through as a graph
 - Marking like mark-and-sweep, but using 3 colours (white, grey, black)
- **Generational garbage collectors**
 - Objects that have been alive for a long time will probably stay for some time more
 - Objects grouped as generations based on creation times

Incremental Collector Tri-Color Marking



Automatic garbage collection in RTSJ

- “Garbage collector is independent and can be changed”
 - RTSJ does not specify any GC, but gives 2 examples of how GC should be implemented
- “Allow the program to precisely characterize an implemented GC algorithm’s effect on the execution time, preemption, and dispatching of real-time Java threads.”
 - GC algorithm should be configurable (scanning rates, CPU usage, priorities ..)

Physical Memory Access

- Embedded applications often require direct memory access for
 - Device drivers
 - Memory-mapped I/O
 - Battery-backed RAM
 - Flash memory
- **RawMemoryAccess** contains methods to create/access a range of physical memory
 - Read-/Write Methods
 - Access based on byte, short, long, float

Synchronization

- Java : *synchronized* keyword
- Communication between NHRT and regular threads needed
- NHRT can not wait for full queues
- Wait free queues
 - Wait-Free-Write-Queue
 - Wait-Free-Read-Queue
 - Wait-Free-Double-Ended-Queue

Priority Inversion

- Default behaviour of synchronized must be: priority inheritance
- RTSJ defines priority ceiling emulation protocol
 - Synchronized segment has a allocated a priority level that indicates the highest possible priority for any thread trying to enter the segment
 - After entering into the segment, the thread's priority is raised to the ceiling value

Handling Posix-Signals

```
public final class POSIXSignalHandler
{
    public static final int SIGABRT;
    public static final int SIGTERM;
    public static final int SIGCANCEL;
    ...
    public static void addHandler(int signal,
                                 AsyncEventHandler handler);
    public static void setHandler(int signal,
                                 AsyncEventHandler handler);
}
```

Realtime Security

- “System and Options”
- Primarily to check physical memory access
- Check if the application is allowed to set the scheduler

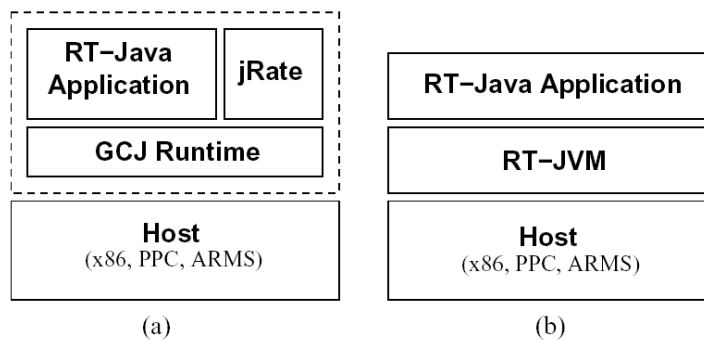
Realtime System

```
public final class RealtimeSystem
{
    public static final byte BIG_ENDIAN
    public static final byte BYTE_ORDER
    public static final byte LITTLE_ENDIAN
    public static GarbageCollector currentGC
    ()
    public static void setSecurityManager
    ( RealtimeSecurity manager)
}
```


RTSJ Implementations

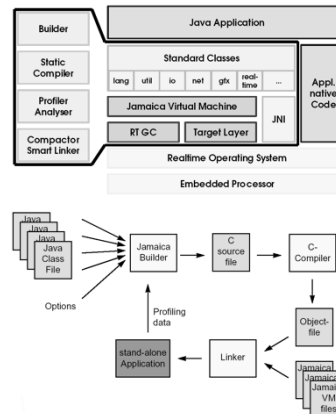
- **Reference Implementation by TimeSys**
 - <http://www.rti.org>
 - Based on Timesys Real-Time Linux / x86
- **Mackinac: Sparc/x86 running Solaris 10**
- **Open source implementation, jRate**
 - <http://tao.doc.wustl.edu/~corsaro/jRate/>
 - PhD thesis of Angelo Corsaro
- **JamaicaVM - aicas GmbH (Karlsruhe)**
- **Esmertec Jbed**
 - 256 kByte including RTOS
- **aJile Systems aj-100**
 - Hardware implementation

jRate – Overview Precompilation



JamaicaVM – aicas GmbH JamaicaVM

- Java bytecode interpreter
- 128 Kbyte minimal footprint
- Real-time garbage collector
- Implements RTSJ exept
 - PriorityCeilingEmulation
 - VTPhysicalMemory
 - LTPhysicalMemory
 - ImmortalPhysicalMemory

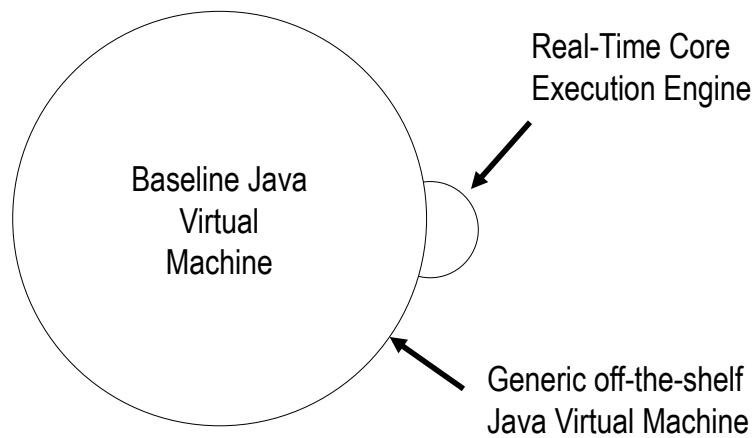


Targets: Sun Solaris, VxWorks, Windows, RTEMS, INTEGRITY

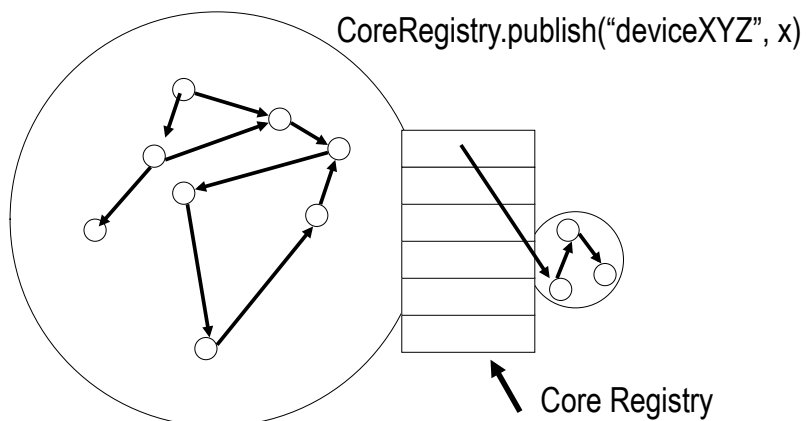
J Consortium

- HP, Aonix, Ericsson, Microsoft, Mitre and NewMonics
- Real-Time Java Working Group
- Core Real-Time Extensions for Java
- Specifies performance like C++
- <http://www.j-consortium.org/rtjwg/index.shtml>

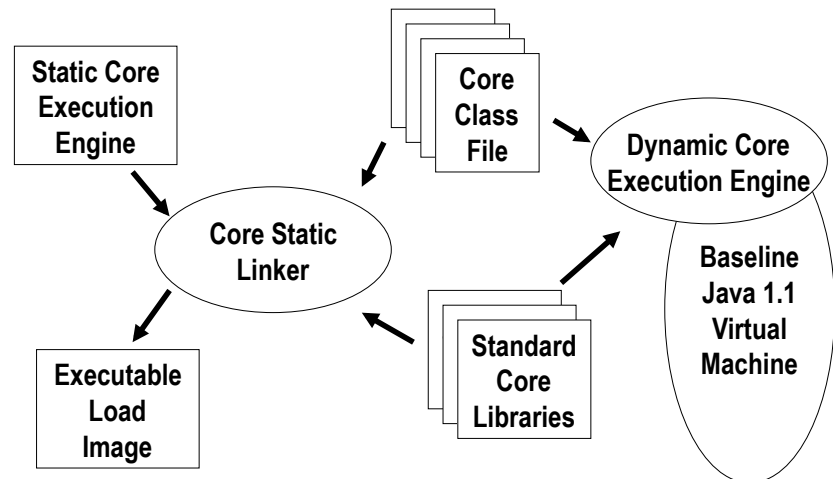
RTJWG View of the Core



The Core Registry



Core Deployment



Distributed RTSJ

- **State : Java Specification Request**
- **RMI for Real-Time application communication**
 - Predictable end-to-end timeliness
 - Other trans-node properties
- **Specification of flow control mechanisms**
- **"The Distributed Specification for Java – An Initial Proposal", E. Douglas Jensen**
- **Following : OMG Dynamic Real-Time CORBA**

Literature

- Real-Time Specification for Java
 - <http://www.rti.org/>
- The Real-Time Java Platform : Mackinac White Paper
- Seminar on Real Time Linux and Java - Spring 2001
 - <http://www.cs.helsinki.fi/u/kraatika/Courses/rt-sem01s.html>
- “Real-Time Java Platform Programming”, Peter Dribble, Prentice Hall PTR
- Sun Java Real-Time System (Java RTS)
 - <http://java.sun.com/javase/technologies/realtime>