Trust Based Access Control (TBAC)

Vincent Schwarzer

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Trust and Reputation

Protect your data with new levels of endpoint security

etwork security schemes center on s perimeter-based solutions. However, the g exposure comes from unsecured endpoints that access the network. Our unique new security and availability solutions help enterprises prevent unauthorized access of corporate data, and can even immediately restore a user's PC in the event of virus attacks or system failures. For more than 25 years Phoenix has been the undisputed leader in BIOS firmware for the PC and computing markets, Now, Phoenix has leveraged its core systems software expertise to create compelling new endpoint security and availability solutions that fill the gaps left by other offerings. To learn how you can achieve new levels of endpoint confidence, please visit www.phoenix.com.



Secure from the START



Cahill, Gray et al. "Using trust for secure collaboration in uncertain environments" (2003) [1]

Trust an Introduction

- Many researchers in fields of psychology, sociology & philosophy have studied the concept of Trust (e.g. Deutsch (1962), Gambetta (1988), Knight and Chervany(1996))
- No unified definition, depends on
 - o authors viewpoint
 - o context
- Trust is:
 - Subjective notion
 - o individual
 - o not symetric
 - situation specific (context needed)
 - \circ self preserving
 - o self amplifying
- Trust is inherently linked to risk

Trust and Reputation: Definition Trust

2 Kinds of Trust:

Reliability Trust (Gambetta (1988))[2]

Trust is the subjective probability by which an individual, A, expects that another individual, B, performs a given action on which its welfare depends.

Decision Trust (McKnight & Chevrany (1996))[3]

Trust is the extent to which one party is willing to depend on something or somebody in a given situation with a feeling of relative security, even though negative consequences are possible.

Trust and Reputation: Definition Reputation

Reputation is what is generally said or believed about a person's or thing's character or standing. (Oxford Dictionary)

Difference Reputation and Trust

- Josang, Ismail, Boyd "A Survey of Trust and Reputation Systems for Online Service Provision" (2007) [4]
- (1) "I trust you because of your good reputation."
- (2) "I trust you despite your bad reputation."
- (1) Uses Public Information to base his trust in the trustee.
- (2) Relying party has some Private Information about the trustee e.g. through direct experience that overrule any reputation that a person might have.

Difference Reputation and Trust

Trust is:

- personal and subjective
- bases on different factors that are different weighted

Reputation is:

- collective measure of trustworthiness
- based on ratings and referrals
- can relate to a group or to an individual

Difference Trust and Reputation Systems

	Trust Systems	Reputation Systems		
Score	Reflects the relyings party subjective view of an entitys trustworthiness	Public Reputation Score as seen by the whole community		
Transitivity	Explicit Component	Implicitly taken into account		
Input	Subjective and general measures of trust	Specific and Objective Events (e.g. Transactions)		

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2. Computing Trust and Reputation

• Physical world trust and reputation systems don't work in IT systems

Reason:

- traditional cues of trust and reputation are missing
- trust and reputation information exchange constrained to local group in physical world

\rightarrow electronic substitutes needed

2. Reputation System Properties

• Proposed by Paul Resnick et.al (2000) [5]

Properties:

- 1. Entities must be long lived, so that with every interaction there is always an expectation of future interactions.
- 2. Ratings about current interactions are captured and distributed.
- 3. Ratings about past interactions must guide decisions about current interactions.

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3. Schema



Source: Mármol; Pérez (2009) [6]

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4. Topologies

Centralised Reputation System	Distributed Reputation System				
 Centralised communication protocols → that participants can provide reputation ratings /obtain rep. scores 	 Distributed communication protocol → Obtain ratings from other particpants 				
 A reputation computation engine → used by central authority to derive reputation scores 	 A reputation computation engine → used by each participant to derive rep. scores of target parties 				



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5. Examples

2 Examples:

- Beta Reputation System (Centralised Reputation System)
- Eigentrust (Distributed Reputation System)

Pre-assumptions:

 Some authentication mechanism is in place to prevent that peers change their identity Example 1.: Selecting Reliable Service Provider in Centralised System with Beta Reputation System

Example 1: Beta Reputation System

- Part of Probability Based Trust Models
- Proposed by Jøsang & Ismail (2002) [7]
- Combines Beta Probability Density Function (PDF) and Belief Model proposed by Jøsang & Audun (2001) [8]
- Based in the theory of statistics
- Universally useable (e.g. E-Commerce Systems for Rating, TBAC)

Example 1: The System



- Reputation Rating Centre has initial seed value (1,1)
- Feedback is given as a pair (r,s) of continous values to express degree of satisfication (r) and dissatisfication (s)
 16.07.13
 Vincent Schwarzer

Example 1: Reputation Rating



Re p(
$$r_A^F, s_A^F$$
)= $\frac{3-1}{3+1}$ = 0,5
Re p(r_B^F, s_B^F)= $\frac{3-4}{3+4}$ =-0,143

Transaction	Α	В	С	D	E	
F	(3,1)	(3,4)	(1,2)	(1,1)	(2,3)	

Example 1: Reputation Discounting

- Part of Jøsang & Audun 2001 proposed Belief Model [8]
- Feedback from highly reputed agents should carry more weight than feedback from agents with low reputation rating
- Metric "Opinion" to describe beliefs about the truth of statements

$$\omega_{\mathbf{X}}^{\mathbf{A}} = (\mathbf{b}, \mathbf{d}, \mathbf{u}) \quad \mathbf{b}, \mathbf{d}, \mathbf{u} \in [0, 1]$$

- b = belief, d = disbelief, u= uncertainty \rightarrow b+d+u = 1
- Possible to map b,d,u to the presentation of r and s

Example 1: Reputation Discounting



Example 1: Forgetting

- Old Feedback not relevant for actual reputation rating
- Old Feedback is given less weight than more recent feedback by using a forgetting factor λ
- $\lambda = 1$ not having forgetting factor; $\lambda = 0$ only last feedback value counts



Example 1: Beta Reputation System

- Reputation Rating
- → Calculate Reputation Value based on Collected Feedback
- Reputation Discounting → Given Feedback has different weight based on collected Feedback in Reputation Centre
- Forgetting

→ Old Feedback discounted by factor λ
 for each new feedback tuple

Example 2: EigenTrust Distributed System to calculate Trust in P2P networks

EigenTrust: General Information

 Eigentrust Algorithm proposed by Kamvar, Schlosser and Garcia Molina (2005) [9,10]

Problem:

Inauthentic Files distributed by malicious peers on a P2P network

<u>Goal:</u>

Identify malicious peers that provide inauthentic files

Method:

Each peer *i* gets a unique global trust value that refletcs the experience of all peers in the network with the peer *i*.

EigenTrust: Terminology

- Normalized Local trust value: c_{ij}
 - Sum of Ratings (postive/negative) of the individual transactions that peer i had with peer j in the past.
- Global trust value: t_i
 - \circ The trust that the entire network places on peer *i*.
- Local Trust Vector \hat{C}_{i}
 - $\circ\,$ Vector with all aggregated and normalized local trust values $S_{ij}^{}$ peer i has about other peers j
- Global Trust vector \vec{t}
 - \circ contains all t

EigenTrust: Calculation

- 1. EigenTrust
 - 1.1. Calculate Local Trust Values
 - 1.2. Normalizing Local Trust Values
 - 1.3. Aggregating Local Trust Values
- 2. Basic EigenTrust
- 3. Distributed EigenTrust
- 4. Secure EigenTrust

1.1. EigenTrust: Calculate Local Trust Values Sij



Peer 4

Peer 3

• local trust value depends on number of positive (tr(i, j) = 1) and negative (tr(i, j) = -1) transactions between peer *i* and *j*

• Local trust calculation:



				- 1		•													
P1	+	-	s_{ij}	P2	+	-	s_{ij}	P3	+	-	s_{ij}	P4	+	-	s_{ij}	P5	+	-	s_{ij}
1				1	1		1	1	2	-2	0	1				1			
2	1		1	2				2	3	-1	2	2	3		3	2	1	-2	-1
3	2		2	3	2		2	3				3	2	-1	1	3			
4				4	2	-1	1	4				4				4	4		4
5				5	3		3	5				5	3	-1	2	5			

1.2. EigenTrust: Normalizing EigenTrust

- Problem: Malicious Peers could assign arbitrarily high local trust values to other malicious peers and arbitrarily low trust values to other peers.
- **Solution:** Normalize the Local Trust Values

$$C_{ij} = \frac{\max(s_{ij}, 0)}{\sum_{j} \max(s_{ij}, 0)}$$
$$C_{i1} + C_{i2} + \dots + C_{in} = 1 \quad \{C_{ij} \in \Re | 1 \le C_{ij} \le 0\}$$

Drawbacks: No distinction between if peer *i* had with peer *j* no or poor interaction

1.2. EigenTrust: Normalizing EigenTrust



Peer 4

Peer 3

Written as vector for Peer 1:

0 1/3 2/3 0 0

 C_{5j}

0

0

0

1

0

				- 1		-												
P1	+	-	C_{1j}	P2	+	-	C_{2j}	P3	+	-	C_{3j}	P4	+	-	C_{4j}	P5	+	-
1			0	1	1		1/7	1	2	-2	0	1			0	1		
2	1		1/3	2			0	2	3	-1	1	2	3		3/6	2	1	-2
3	2		2/3	3	2		2/7	3			0	3	2	-1	1/6	3		
4			0	4	2	-1	1/7	4			0	4			0	4	4	
5			0	5	3		3/7	5			0	5	3	-1	2/6	5		

1.3. EigenTrust: Aggregating EigenTrust

- Each peer bases its choice of downloads using its own opinion vector *ci*
- Depending on the previous experience with the other peer determines the chance how likely peer i will choose peer j for further transactions

Problem:

Each peer has only limited past experience and knows only few other peers.

Solution:

Ask other peers **i** trust about their opinion of other peers **j** and weight their opinion .





1.3. EigenTrust: Know all Peers

<u>Know all Peers:</u> Ask your friends: $t_i = C^t c_i$

 $\frac{\text{Ask friends' friends:}}{t_i = (C^T)^2 c_i}$

<u>Keep asking until t converges:</u> $t_i = (C^T)^n c_i$

- for large **n** all t_i converges to t: $\vec{t} = \vec{t}_i, \forall i$
- peers can cooperate to compute and store t

2. Basic EigenTrust

A priori notion of trust

- Some peers in the network are known to be trustworthy (e.g., founders)
- \circ Define distribution \vec{p} over pre-trusted peers
- $\circ~$ Use $\vec{p}~$ as initial vector

Used for:

Inactive peers

- o If a peer doesn't interact or has only bad interactions
- Trusts always Pre Trusted peer

Malicious collectives

 Group of peers who knows each other, who give each other high local trust values and give all other peers low local trust values in an attempt to subvert the system

$$\vec{t}^{(k+1)} = (1-a)C^{T}\vec{t}^{(k)} + a\vec{p}$$

3. EigenTrust: Distributed EigenTrust

 A_i : set of peers which have downloaded files from peer i B_i : set of peers from which peer i has downloaded files

Distributed EigenTrust Algorithm

for each peer i ask peer j \in A_i for cji and t_j(0) = p_j **repeat Compute** $t_i^{(k+1)} = (1-a)(c_{1i}t_i^{(k)} + ... + c_{ni}t_n^{(k)}) + ap_i$ **send** your opinion cij and trust value $t_i^{(k+1)}$ to all peers $j \in B_i$ **wait** for all peers $j \in A_i$, to respond with their opinion c_{ji} and trust value $t_j^{(k+1)}$ **until** $|t(k+1) - t(k)| < \epsilon$

4. EigenTrust: Secure EigenTrust

Problem: Malicious peer can cooperate and/or report false value to other peers

- **<u>Goal</u>**: Peer should not hold its own t_i
 - t_i should not be computed by only one peer

Solution:

- Multiple score managers computes combined a peers global trust value
- verified by comparing the calculated results and choose the majority of the results (Distributed Hash Table)

EigenTrust: Choose a Peer

- Choose a peer with a probability corresponding to its global trust value
- With a probability of 10% choose a peer with a global trust value of zero.

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6. Problems / Attacks in Reputation System

- Mármol, Pérez "Security threats scenarios in trust and reputation models for distributed systems" (2009) [6]
- Low Incentive for Providing Rating
- Bias Towards Positive Rating
- Unfair Ratings
- Change of Identities
- Quality Variations over Time
- Ballot Box Stuffing
- Malicious Peers/Collectives (Malicious Collective, Campuflaged Collective,...)

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Why is TBAC necessary?

- Liu; Trust-Based Access Control for Collaborative System; 2008 [11]
- Existing access control models (DAC,MAC,RBAC) well suited for centralized/static environments
- These models are not suitable for collaborative enviroments
 - members and ressources are dynamical
 - o too high maintenance effort for human being

Conclusion:

Use a model of human notion of trust and community as the basis of assigning privileges. Rights/Privileges are dynamically assigned based on Risk of an action (context) for peers/ressources in a network.

Realisation:

Trust Based Access Control

TBAC in General

- No widely applied TBAC Standards yet:
 SECURE Project
- 2 approaches for implementing TBAC currently:
 - Trust through certifcate based system
 - Trust Computation based on Transaction Ratings
- Most of the current approaches extend the Role Based Access Control (RBAC) with the Notion of Trust

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Examples for TBAC

- Stackoverflow (Meta Stackoverflow)
- TBAC System for File Access



Example 1:

- Users get Reputation for different acitivities in the community
- Own Sub-site to discuss the system called Meta stackoverflow
- Centralised Reputation System using Simple Summation of Ratings
- Incentive to be active on the website and used as filtering mechanism for malicious/missbehaving users

Action (Created by User)	Rep.	Action (Created by User)	Rep.
Question Voted Up/Useful	+ 5	Question or Answer is voted <i>down/</i>	- 2
Answers is voted Up/Useful:	+ 10	not userui	
One of your Answers become	+ 15	You vote an answer <i>down/not</i> useful	- 1
Remove a downvote from an Answer	+ 1	Upvote on one of your questions is removed	- 5



Calculation

$$Reputation = \sum_{i}^{n} Rep(Action)$$

- Thresholds defined for different Privileges
- If certain amount of Reputation is reached additional Privileges are granted

	Privilege	Points Required					
	Create Posts / Comment Everywhere	1					
	Participate in per-site meta	5					
	Remove new user restrictions	10					
6.07	1 Frusted User Vincent S	chwarzer 20 000					

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Example 2: TBAC System for File Access

 Feng,Lin,Peng,Li ;"A Trust and Context Based Access Control Model for Distributed Systems" (2008) [13]



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Problems and Open Questions

Trust Computation Systems

- How to get the initial trust values?
- No established de-facto Standards
- How to choose Pre-Trusted Peers

Research Field

- Many "low quality" Research Papers
- Most researchers develop their own systems further no crossover or collaboration between the systems

Thank you for your attention!

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